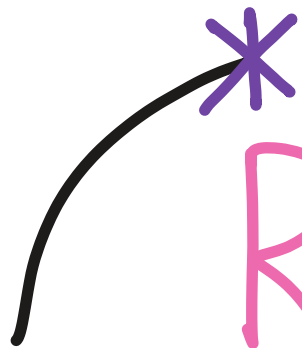




Relational Cryptis

A Separation Logic for Indistinguishability
in the Symbolic Model

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A motivating example

Alice  generates a key k

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 gives Eve  oracle access to $\text{Enc}_k(\cdot)$

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 tosses a coin $\textcircled{1}$

 sends $\text{Enc}_k(M_{\textcircled{1}})$ to 

A motivating example (IND-CPA)

A cryptosystem is indistinguishable under chosen-plaintext attack if  cannot determine whether $Enc_k(M_0)$ or $Enc_k(M_1)$ was sent.

The Dolev-Yao Model

— Terms, not bitstrings

Inductive term :=

| TInt of Z

| TPair of term & term

| TKey of key-type & term

| TSeal of term & term

⋮

The Dolev-Yao Model

- Terms, not bitstrings
- Idealized crypto using rewrite systems

$\text{adec } k \text{ (aenc pkey } k) m = \text{Some } m$

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- Attacker with unlimited computational power
 - + receive, send, duplicate, manipulate messages sent over the network

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- Idealized crypto using rewrite systems
- Attacker with unlimited computational power
 - + receive, send, duplicate, manipulate messages sent over the network
 - + bound by the rewrite rules

Cryphis

— Separation logic for symbolic cryptography

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Cryphis

- Separation logic for symbolic cryptography
- $\text{public } t \multimap^* \text{WP send } c \ t \ \{\{\psi\}\}$

Definition $\text{public } t : \text{iProp} :=$

match t with

| $\text{TInt } t \Rightarrow \text{True}$

| $\text{TKey } kt \ t \Rightarrow \text{public_key_type } kt$

| $\text{TSeal } (\text{TKey } \text{AEnc } t) \ t' \Rightarrow$

$\square (\text{public } (\text{TKey } \text{ADec } t) \Rightarrow \text{public } t')$

:

Cryphis

- Separation logic for symbolic cryptography
- $\text{public } t \multimap^* \text{WP send } c \ t \ \{\{\psi\}\}$
- The public predicate enforces rewrite rules
$$\begin{aligned} \lceil \text{adec } k \ t = \text{Some } t' \rceil &\multimap^* \text{public } k \multimap^* \text{public } t \\ &\multimap^* \text{public } t' \end{aligned}$$

Cryphis

- Separation logic for symbolic cryptography
- public $t \rightarrow^* \text{WP send } c \ t \ \{\{\psi\}\}$
- The public predicate enforces rewrite rules
 - Can prove term secrecy, agreement etc.

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Cryphis

- Separation logic for symbolic cryptography
- public $t \rightarrow^* \text{WP send } c \ t \ \{\{\psi\}\}$
- The public predicate enforces rewrite rules
 - Can prove trace properties
 - Cannot prove hyperproperties



IND-CPA?

— What does IND-CPA even mean in our setting?



IND-CPA?

- What does IND-CPA even mean in our setting?
- Observational equivalence?



ReLoC

- Separation logic on top of Iris for program equivalence.



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- $e \lesssim_{\text{ctx}} e' : \tau$



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- $e \lesssim_{\text{ctx}} e' : \tau$
- $\text{REL } e \ll e' : \tau \rightarrow e \lesssim_{\text{ctx}} e' : \tau$



ReLoC

- Separation logic on top of Iris for program equivalence.
- $e \lesssim_{\text{ctx}} e' : \tau$
- $\text{REL } e \ll e' : \tau \rightarrow e \lesssim_{\text{ctx}} e' : \tau$
- Attacker knowledge? τ ?

Composition

Lemma foo : REL $e \ll e' : \tau$



Lemma foo' : $(\tau \rightarrow \psi) \rightarrow$
REL $e \ll e' : \psi$

Attacker knowledge

- Bijection between terms exposed to the network.

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- Must preserve Dolev-Yao properties

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- *public* as a bijection?

Attacker knowledge

- Bijection between terms exposed to the network.
- Must preserve Dolev-Yao properties
- public as a bijection?

 $\text{TSeal}(\text{TKey AEnc } k) t, \text{TSeal}(\text{TKey AEnc } k') t'$

 $\text{TSeal}(\text{TKey AEnc } k) t_1, \text{TSeal}(\text{TKey AEnc } k'_1) t'_1$

 $\text{TKey ADec } k, \text{TKey ADec } k'$

publicly-related

$\text{PUB} \langle t, t' \rangle \dashv\vdash \text{match } t, t' \text{ with}$

| TInt $n, \text{TInt } n' \Rightarrow \ulcorner n = n' \urcorner$

| TPair $t_1 \ t_2, \text{TPair } t_1' \ t_2' \Rightarrow \text{PUB} \langle t_1, t_1' \rangle \wedge$
 $\text{PUB} \langle t_2, t_2' \rangle$

| TNonce $a, \text{TNonce } a' \Rightarrow$

$\text{public_rel} \hookrightarrow_{\text{BIT}} \square \langle t, t' \rangle$

.

Ghost state 

public-rel $\hookrightarrow_{BIT} \square \langle t, t' \rangle$

— partial bijection

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public-rel $\hookrightarrow_{BIT}$ $\square$ $\langle t, t' \rangle$

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- elements track terms related by fiat

Ghost state

public-rel $\hookrightarrow_{BIT}$ $\square$ $\langle t, t' \rangle$

- partial bijection
- elements track terms related by fiat
- publicly-related does the rest by structural recursion

publicly-related

$\text{PUB}\langle t, t' \rangle \dashv\vdash \text{match } t, t' \text{ with}$

$\vdots$

$\mid \text{TSeal } k \ t_1, \text{TSeal } k' \ t'_1 \Rightarrow$
 $(\text{PUB}\langle k, k' \rangle \wedge \text{PUB}\langle t_1, t'_1 \rangle) \vee$

$\left(\text{public_rel } G_{\text{BIJ}} \square \langle t, t' \rangle \wedge \right.$
 $\text{PRIV}\langle k, k' \rangle \wedge \text{PRIV}\langle t_1, t'_1 \rangle \wedge$
 $\square (\text{match } k, k' \text{ with } \dots$
 $\mid \text{TKey AEnc } k_1, \text{TKey AEnc } k'_1 \Rightarrow$
 $\text{PUB}\langle k_1, k'_1 \rangle \rightarrow \text{PUB}\langle t_1, t'_1 \rangle) \left. \right)$

privately-related

$\text{PRIV} \langle t, t' \rangle$

"t and t' shall be united once more"

— $\neg \neg$ Sybill Trelawney

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$\forall t'', \text{PUB} \langle t, t'' \rangle \rightarrow \lceil t' = t'' \rceil$

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— decreasing measure : $\min(\text{tsize } t, \text{tsize } t')$

$\forall t'', \text{PUB} \langle t, t'' \rangle \rightarrow \lceil t' = t'' \rceil$ 

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$\forall t'', \text{PUB} \langle t, t'' \rangle \rightarrow \lceil t' = t'' \rceil$ ~~X~~

$\forall t'', \triangleright \text{PUB} \langle t, t'' \rangle \rightarrow * \triangleright \lceil t' = t'' \rceil$

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$\forall t'', \triangleright \text{PUB} \langle t, t'' \rangle \rightarrow * \triangleright \lceil t' = t'' \rceil$ 



IND-CPA

Lemma $\text{rel_alice } c \ c' \ b :$

$\text{cryptis_rel_ctx} \rightarrow *$

$\text{channel_rel } c \ c' \rightarrow *$

$\text{REL_alice } c \ \ll \ \text{alice } c' :$

$\lambda (\text{choice}, \text{guess}) (\text{choice}', \text{guess}')$

$b = \text{choice} \rightarrow b \neq \text{choice}' \wedge \text{guess} = \text{guess}'$



IND-CPA

— create sk_A, sk_A'



IND-CPA

- create skA, skA'
- send $PUB\langle pkey\ skA, pkey\ skA' \rangle$
+ relate by fiat, not structure



IND-CPA

- create $sk_A, sk_{A'}$
- send $\text{PUB}\langle \text{pkey } sk_A, \text{pkey } sk_{A'} \rangle$
- receive $\text{PUB}\langle m_0, m_0' \rangle, \text{PUB}\langle m_1, m_1' \rangle$



IND-CPA

- create $sk_A, sk_{A'}$
- send $PUB\langle pkey\ sk_A, pkey\ sk_{A'} \rangle$
- receive $PUB\langle m_0, m_0' \rangle, PUB\langle m_1, m_1' \rangle$
- toss coin b



IND-CPA

- create $sk_A, sk_{A'}$
- send $PUB\langle pkey\ sk_A, pkey\ sk_{A'} \rangle$
- receive $PUB\langle m_0, m_0' \rangle, PUB\langle m_1, m_1' \rangle$
- toss coin b
- send $PUB\langle aenc\ (pkey\ sk_A)\ m_b,$
 $aenc\ (pkey\ sk_{A'})\ m_{\bar{b}} \rangle$



IND-CPA

$\text{PUB} \langle M_0, M_0' \rangle$
 $\text{PUB} \langle M_1, M_1' \rangle$

- send $\text{PUB} \langle \text{pkey } sk_A, \text{pkey } sk_{A'} \rangle$
- send $\text{PUB} \langle \text{aenc}(\text{pkey } sk_A) M_b, \text{aenc}(\text{pkey } sk_{A'}) M_{\bar{b}} \rangle$

+ must relate by fiat, not structure



IND-CPA

$\text{PUB} \langle M_0, M_0' \rangle$
 $\text{PUB} \langle M_1, M_1' \rangle$

- send $\text{PUB} \langle \text{pkey } sk_A, \text{pkey } sk_{A'} \rangle$
- send $\text{PUB} \langle \text{aenc}(\text{pkey } sk_A) M_b, \text{aenc}(\text{pkey } sk_{A'}) M_{\bar{b}} \rangle$

+ must relate by fiat, not structure

+ need to prove

$\text{PRIV} \langle \text{pkey } sk_A, \text{pkey } sk_{A'} \rangle$

Problem

$\text{PUB} \langle \text{pkey } \text{skA}, \text{pkey } \text{skA}' \rangle \rightarrow *$

$\text{PRIV} \langle \text{pkey } \text{skA}, \text{pkey } \text{skA}' \rangle$

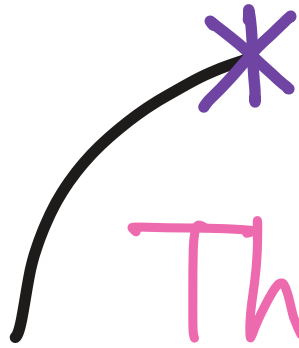
is unprovable!

(expands to a proof of $\Delta P \rightarrow * P$)

Potential Solutions?

- a different definition of PRIV?
- spookier ghost state?

Reach out if you know 🙄



Thank You! Questions?